

ENGR (XMUT) 101

Engineering Technology

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*Te Whare Wānanga
o te Ūpoko o te Ika a Māui*



CAPITAL CITY UNIVERSITY

Week 11 Lecture 1

- Main topics
 - Number system – part 3

Finite Number Representation

- Machines that use 2's complement arithmetic can represent integers in the range:

$$-2^{n-1} \leq N \leq 2^{n-1}-1$$

where n is the number of bits available for representing N .

- Note that $2^{n-1}-1 = (011\dots11)_2$ and

$$-2^{n-1} = (100\dots00)_2$$

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- For 2's complement more negative numbers than positive.
- For 1's complement two representations for zero.
- For a n bit number in base (radix) z there are z^n different unsigned values.

$$(0, 1, \dots, z^{n-1})$$

1's Complement Addition

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Step 1: Add binary numbers

Step 2: Add carry to low-order bit

$$\begin{array}{r} 01100 \\ + 00001 \\ \hline 01101 \\ 0 \\ \hline 01101 \end{array}$$

1's Complement Subtraction

- Subtracting 1's complement numbers is also easy.

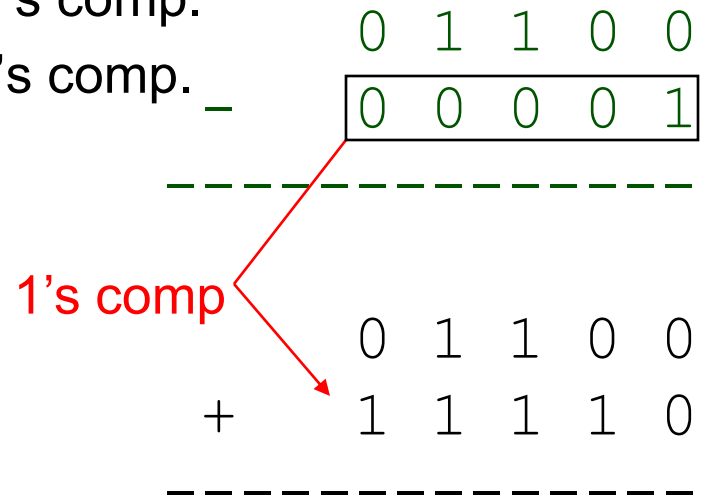
1's Complement Subtraction

- Subtracting 1's complement numbers is also easy.
- Let's compute $(12)_{10} - (1)_{10}$.
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$$\begin{array}{r}
 01100 \\
 - 11110 \\
 \hline
 \end{array}$$

Step 1: Take 1's complement of 2nd operand.

Add

$$\begin{array}{r}
 01100 \\
 + 11110 \\
 \hline
 \end{array}$$

Step 2: Add the binary numbers.

$$\begin{array}{r}
 101010
 \end{array}$$

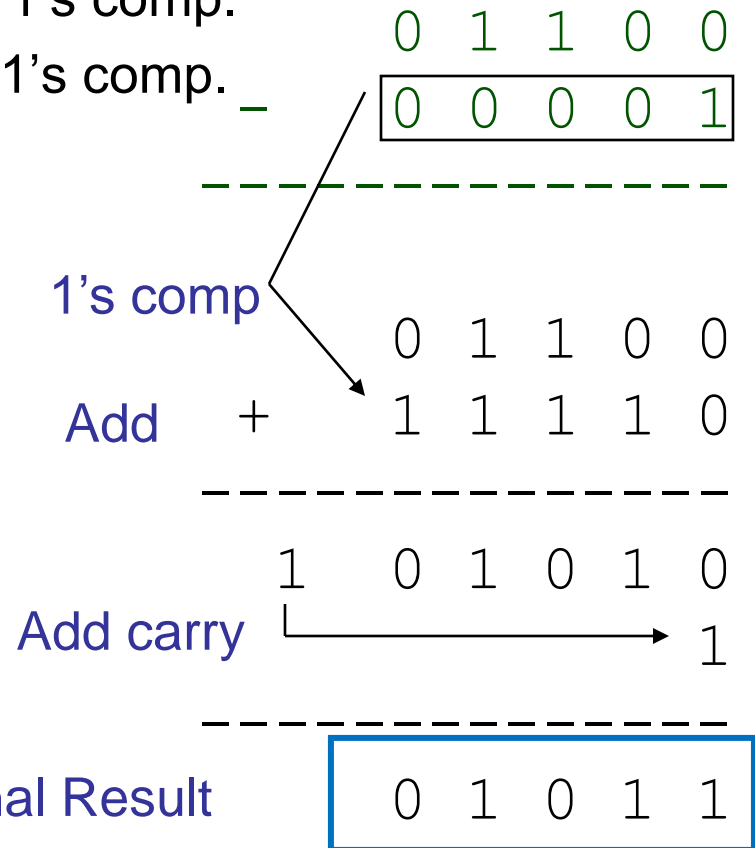
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Step 1: Take 1's complement of 2nd operand

Step 2: Add binary numbers

Step 3: Add carry to low order bit



2's Complement Addition

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2's Complement Addition

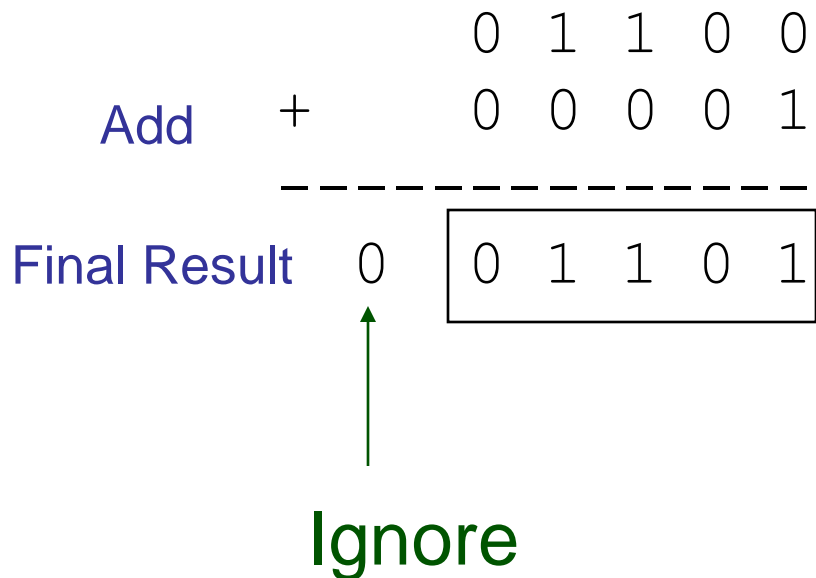
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2's Complement Addition

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- Let's compute $(12)_{10} + (1)_{10}$.
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Step 1: Add binary numbers

Step 2: Ignore carry bit



2's Complement Subtraction

- Follow the 3 steps for subtraction.
- Let's compute $(12)_{10} - (1)_{10}$.
 - $(12)_{10} = +(1100)_2 = 01100_2$ in 2's comp.
 - $(-1)_{10} = -(0001)_2 = 11111_2$ in 2's comp.

2's Complement Subtraction: **Exercise 1**

Compute $(13)_{10} - (5)_{10}$ using the 2s complement form.

5 minutes
to complete this exercise!!

2's Complement Subtraction: **Exercise 2**

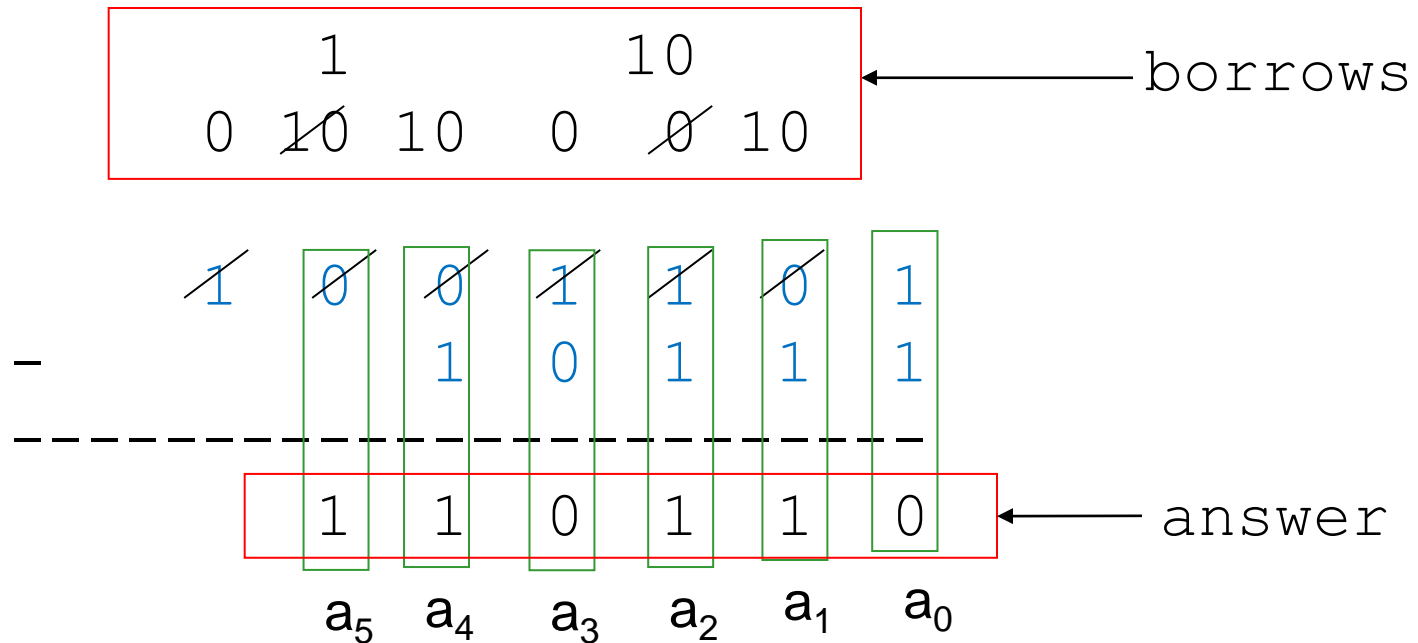
Compute $(5)_{10} - (12)_{10}$

5 minutes

to complete this exercise!!

Binary Subtraction

- We can also perform subtraction (with borrows in place of carries).
- Let's subtract $(10111)_2$ from $(1001101)_2$



Binary Multiplication

What is the decimal equivalent of the 2 binary numbers?

23
X10
230

$$\begin{array}{r} 10111 \\ \times 1010 \\ \hline 00000 \\ 10111 \\ 00000 \\ 10111 \\ \hline 11100110 \end{array}$$

Binary Multiplication Exercise 5.1

Multiply the 2 binary numbers:

$(1001)_2$ from $(111)_2$

$$\begin{array}{r} \\ X \\ \hline \end{array}$$

2 minutes to complete this exercise!

Binary Multiplication Exercise

Multiply the 2 binary numbers:

$(1001)_2$ from $(111)_2$

$$\begin{array}{r} \\ 1001 \\ 111 \\ \hline 1001 \\ 1001 \\ 1001 \\ \hline 111111 \end{array}$$

Diagram illustrating the binary multiplication of $(1001)_2$ by $(111)_2$. The multiplicand 1001 is multiplied by each bit of the multiplier 111 . The resulting partial products are shown, with the first row of each product labeled "1st row x 1".

